

Security in RINA

IRATI Workshop

Barcelona, Spain

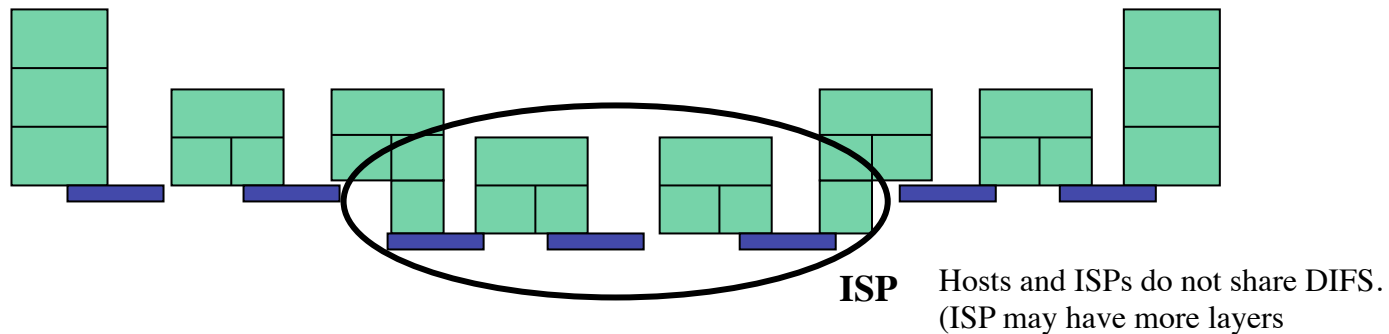
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First a Word on Method

- When trying to work out the IPC Model absolutely no thought was given to security. All of the focus was just understanding the structure.
- People kept asking, What about Security? Is there a security layer?
- Didn't Know. Hadn't thought about it.
- There was the obvious:
 - The recursion of the layer provided Isolation.
 - That only the Application Name and local port-id were exposed to the correspondents.
- Interesting, but hardly an answer
- But it wasn't the time for those questions . . .
- At least not yet . . .

The Recursion Provided Isolation

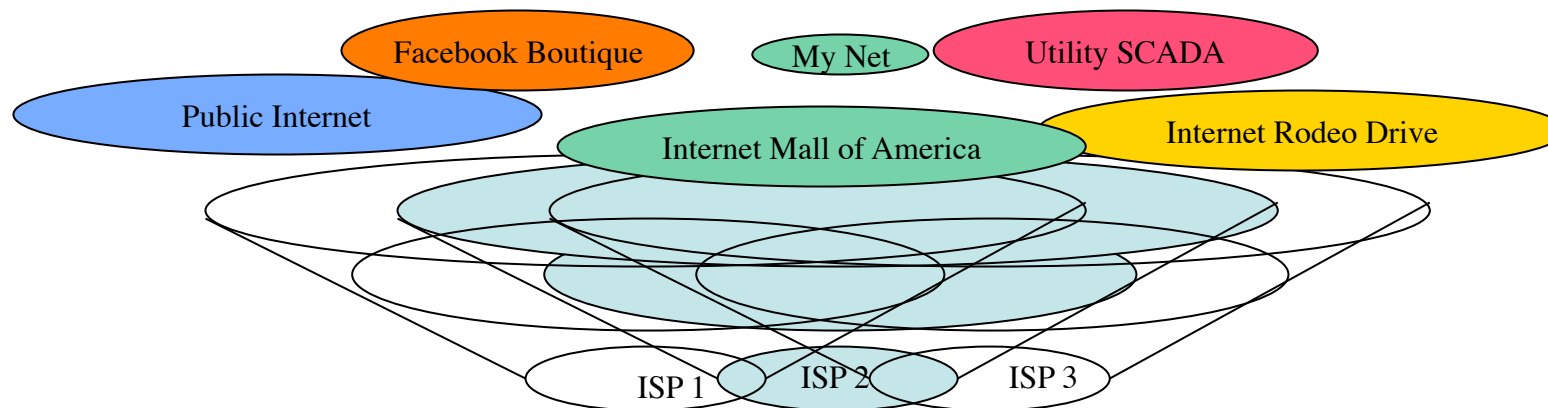


- Security by isolation, (not obscurity)
- Hosts can not address any element of the ISP.
- No user hacker can compromise ISP assets.
 - Unless ISP is physically compromised.

How Does It Work?

Security

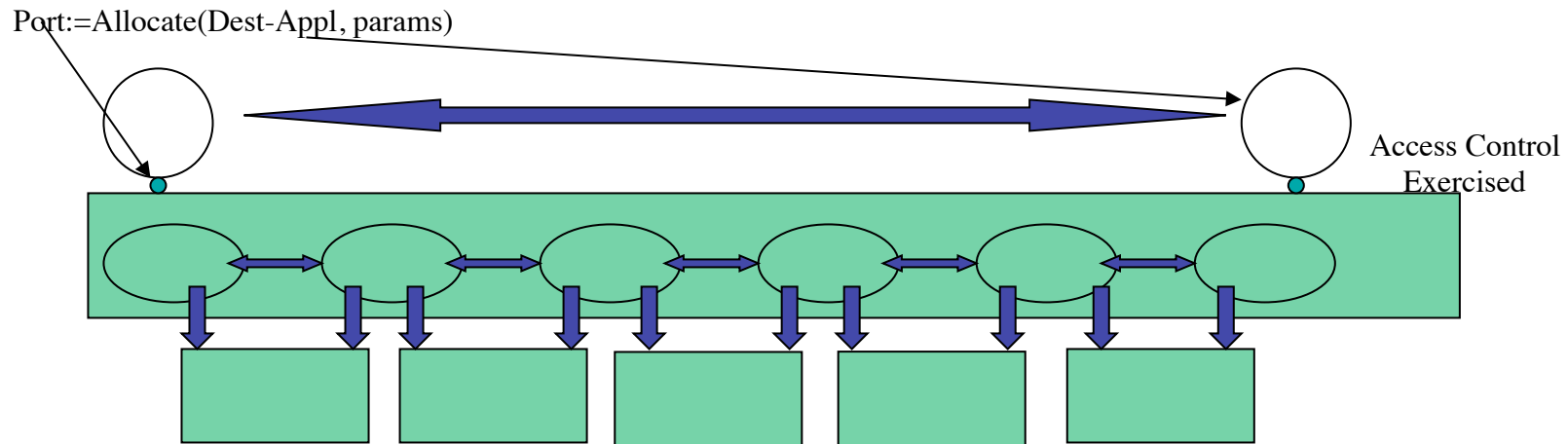
- A Hacker in the Public Internet cannot connect to an Application in another DIF without either joining the DIF, or creating a new DIF spanning both. Either requires authentication and access control.
 - Non-IPC applications that can access two DIFs are a potential security problem.
- Certainly promising



But When It Was Time

- The question was not, how to put in security?
- The question was,
- What does the IPC Model tell us about security?
 - Remember, our first task is always *understanding*.
- Let the Problem Answer the Question!
 - Let the Problem Tell Us What to Do.

The Problem Had a Lot to Say



- We Already Mentioned How Little is Exposed the Layer Above.
- The Original OS Model indicated where Access Control went.
- Creating the Application Connection for Enrollment indicated where Authentication belonged, and that
 - Authentication of Applications must be done by the Applications themselves.
 - All members of the layer are authenticated within policy.
- SDU Protection clearly provided Confidentiality and Integrity.
- That implied that only Minimal trust was necessary:
 - Only that the lower layer will deliver something to someone.

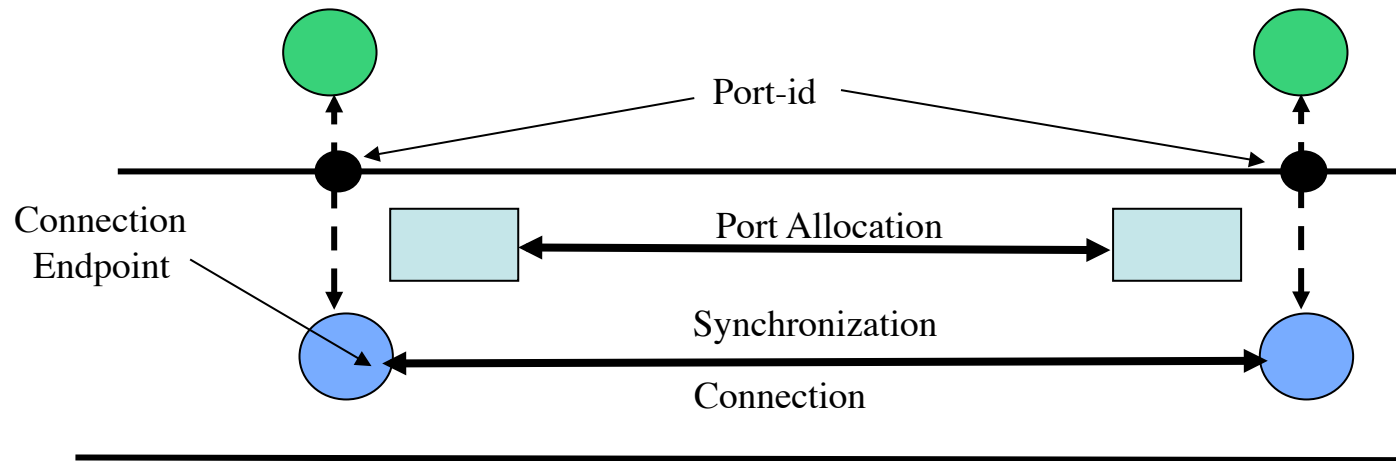
A Very Unexpected Result

- A DIF with no explicit security mechanisms is inherently more secure than the current Internet under the same conditions!
- It would appear that
 - A DIF is a Securable Container.

Other Things Fall Into Place

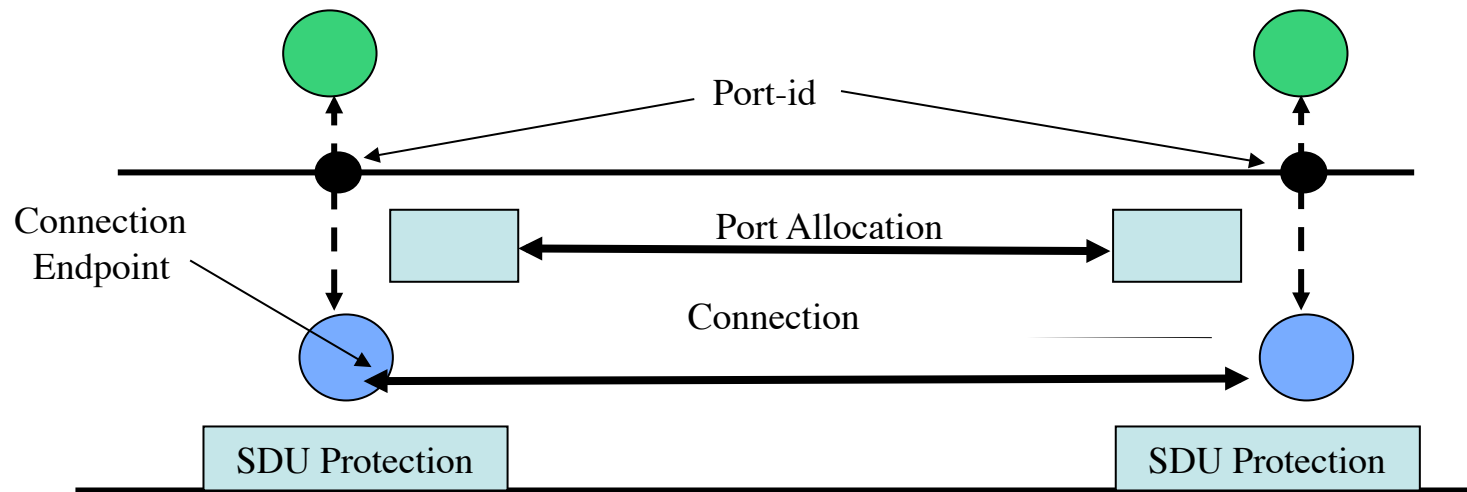
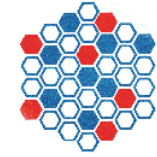
- Data Transfer in RINA is based on Delta-t (Watson, 1980)
- Lot has happened in 30 years, many attacks on TCP have been found:
 - Port scanning
 - SYN attacks
 - Reset Attacks
 - Reassembly Attacks
- Long after delta-t was designed, what about delta-t?
- Short answer:
 - None of them work (Boddapati, et al., 2012)
 - Amazing, totally unexpected
 - Why not?
- Multiple fundamental reasons, but all inherent in the structure:
 - First, have to join the DIF (all members are authenticated)
 - Second, No Well-Known Ports
 - Would have to scan all possible application names!
 - Third and more importantly, . . .

Decoupling Port Allocation and Synchronization



- No Way to Know What CEP-ids are Being Used, Since There is No Relation Between Port-id and CEP-id.
 - Syn Attack: must guess which of 2^{16} CEP-id.
 - Data Transfer: must guess CEP-id and seq num within window!
 - Reassembly attack: Reassembly only done once.

Decoupling Port Allocation and Synchronization: No IPsec

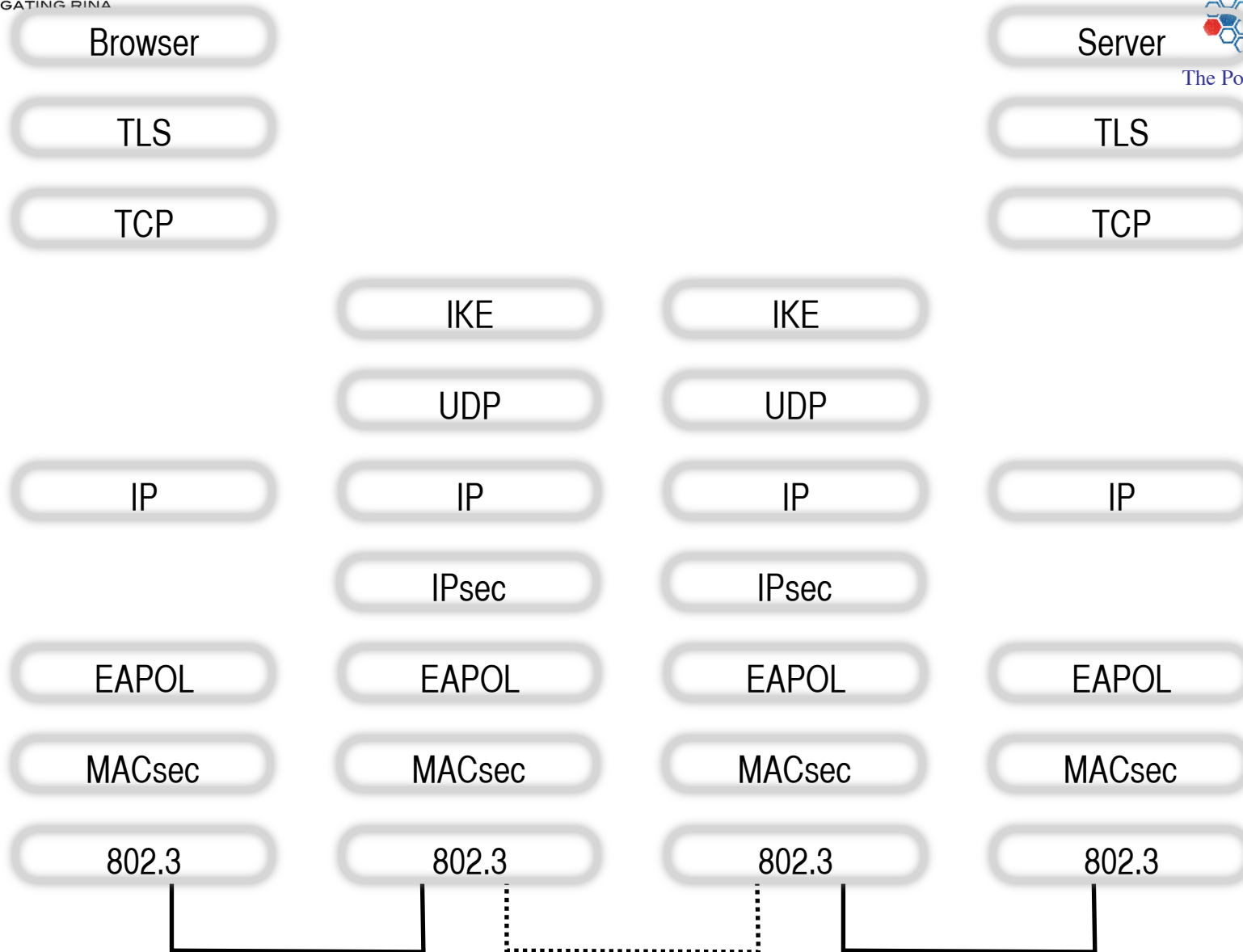


- IPsec is necessary with TCP/IP because no authentication and Sequence numbers turn over too quickly: don't repeat sequence number with same CEP-id.
- With RINA and delta-t, IPC Processes all authenticated, SDU Protection does the encryption, and packet sequence numbers slows rollover, but if it does, then simply allocate a new connection
- And bind it to the same port-ids, old one disappears after 2MPL.

RINA is Inherently More Secure and Less Work



- A DIF is a Securable Container. (Small, 2011)
 - What info required to mount an attack, How to get the info
 - Small does a threat analysis at the architecture level
- Implies that Firewalls are Unnecessary,
 - The DIF *is* the Firewall!
- RINA Security is considerably Less Complex than the Current Internet Security (Small, 2012)
 - Only do a rough estimate counting protocols and mechanisms.
 - See paper for details.



Protocols: 15

Non-Security: 89

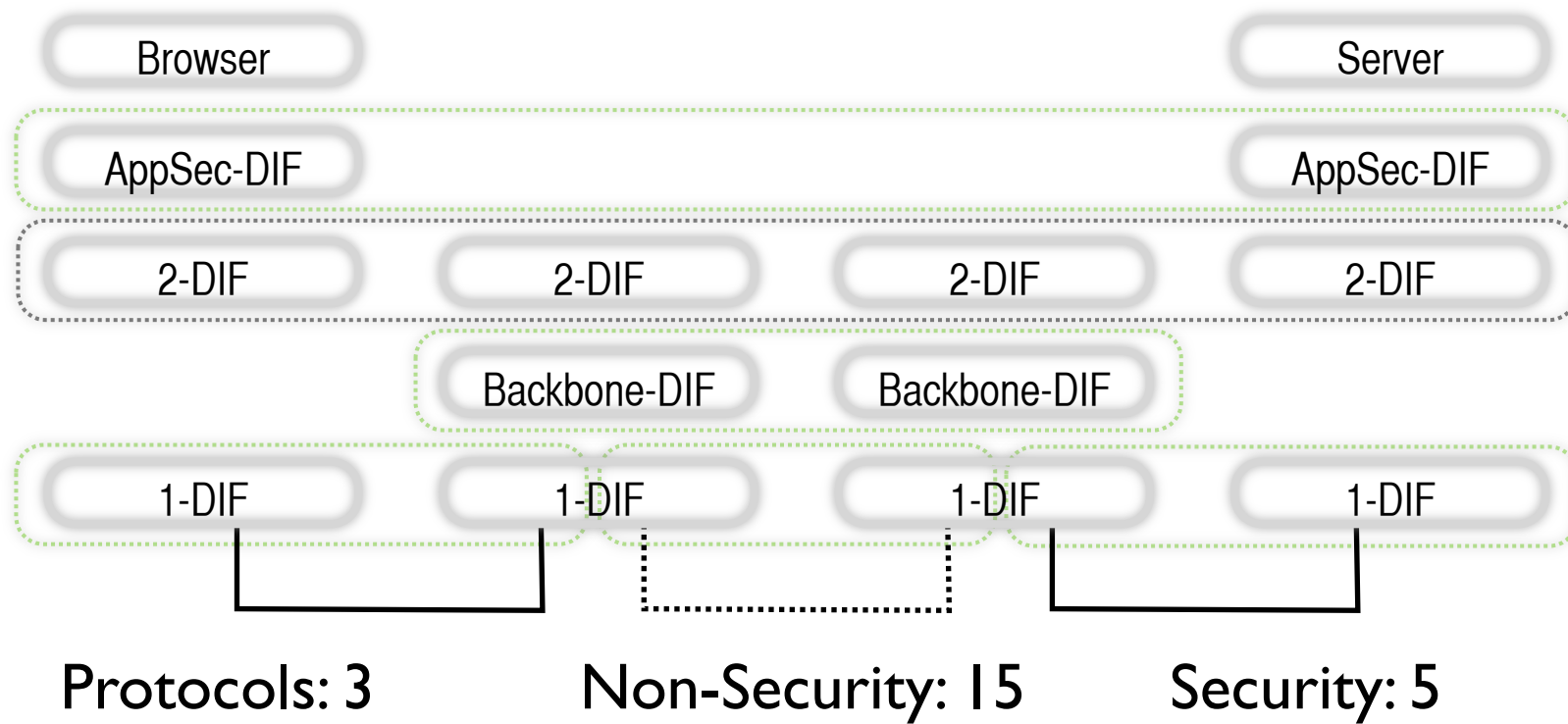
Security: 28

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What Does This Mean?

- Protocols – We Know What That Refers To
- Security Mechanisms – Authentication, Access Control, Integrity, Confidentiality, Non-Repudiation.
- Non-Security Mechanisms – All the others listed in the book: delimiting, relaying, ordering, multiplexing, fragmentation/reassembly, Lost and Duplicate Detection, Flow Control, Retransmission Control, Compression, Addressing, Initial State Synchronization.



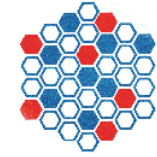
Totals	Internet	RINA
Protocols	15	3
Non-Security Mechanisms	89	15
Security Mechanisms	28	7

To Add Security	Internet	RINA
Protocols	8	0
Non-Security Mechanisms	59	0
Security Mechanisms	28	7

Why Is Internet Security So Bad?

- The Standard Rationale One Sees is that They Didn't Think About It at the Beginning.
 - Neither did We.
 - Nor did Watson.
 - But RINA and delta-t are more secure.
- That Seems to Imply that
 - Good Design May be More Important to Security than Security Is.

Conclusion



- This is a MAJOR Improvement in Internet Security.
 - Not only more secure, but for less cost, with less overhead.
- So is Internet Security solved?
 - Hardly.
 - Still need: to develop the plug-in policy modules
 - to consider DDoS (we have some ideas)
 - As well as protecting against Rogue IPC Processes
 - and much more to explore.
- Most attacks are in the Applications, this does nothing about that.
 - But Much of this applies equally well to DAFs
 - Model implies that OS security reduces to Bounds Checking on Memory and IPC Security.
 - May also make it harder, might be able to deflect more DDoS attacks

Questions?